Windowing Functions

Hitoshi Harada (umi.tanuki@gmail.com) Head of Engineering Dept. FORCIA, Inc. **David Fetter** (david.fetter@pgexperts.com)



What are Windowing Functions?

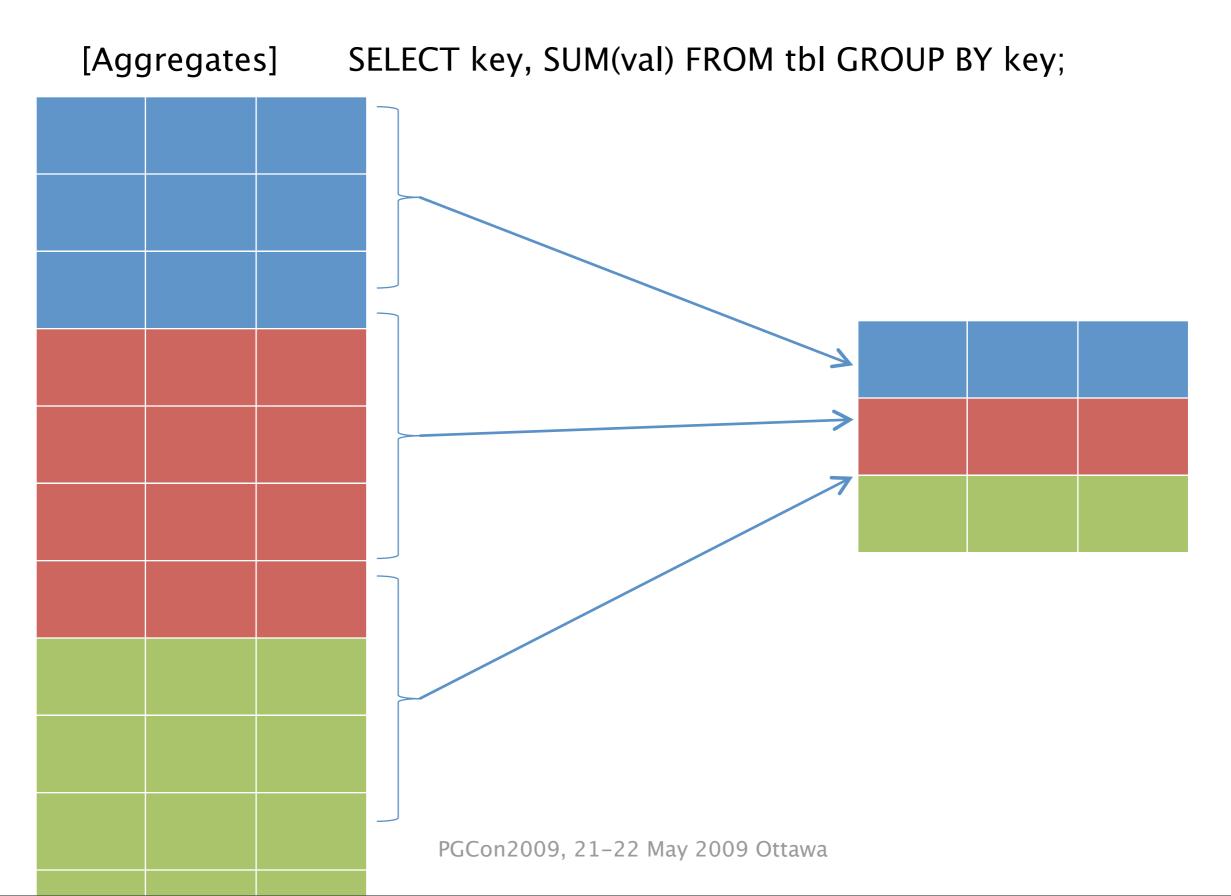
What are the Windowing Functions?

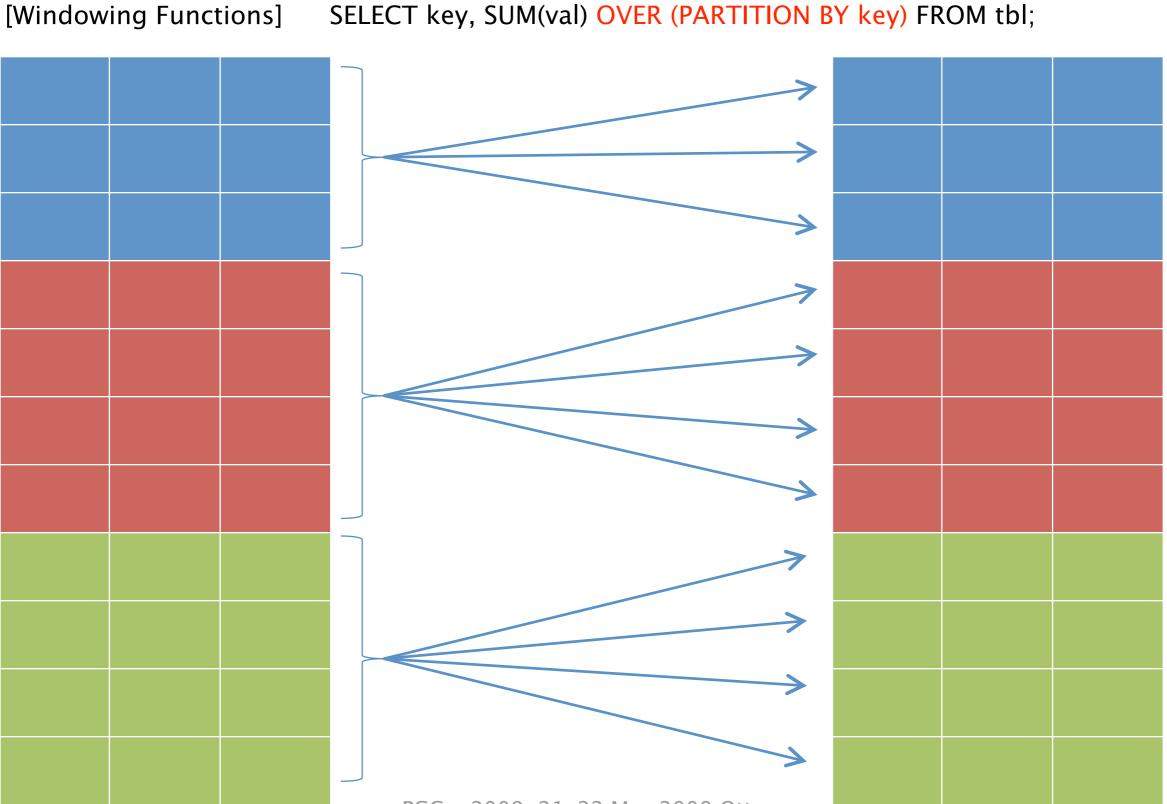
- Similar to classical aggregates but does more!
- Provides access to set of rows from the current row
- Introduced SQL:2003 and more detail in SQL:2008
 - SQL:2008 (http://wiscorp.com/sql200n.zip) 02: SQL/Foundation
 - 4.14.9 Window tables
 - 4.15.3 Window functions
 - 6.10 <window function>
 - 7.8 <window clause>
- Supported by Oracle, SQL Server, Sybase and DB2
 - No open source RDBMS so far except PostgreSQL (Firebird trying)
- Used in OLAP mainly but also useful in OLTP
 - Analysis and reporting by rankings, cumulative aggregates

- Windowed table
 - Operates on a windowed table
 - Returns a value for each row
 - Returned value is calculated from the rows in the window

- You can use...
 - New window functions
 - Existing aggregate functions
 - User-defined window functions
 - User-defined aggregate functions

- Completely new concept!
 - With Windowing Functions, you can reach outside the current row





Who is the highest paid relatively compared with the department average?

depname	empno	salary
develop	10	5200
sales	1	5000
personnel	5	3500
sales	4	4800
sales	6	550
personnel	2	3900
develop	7	4200
develop	9	4500
sales	3	4800
develop	8	6000
develop	11	5200

SELECT depname, empno, salary, avg(salary) OVER (PARTITION BY depname)::int, salary – avg(salary) OVER (PARTITION BY depname)::int AS diff FROM empsalary ORDER BY diff DESC

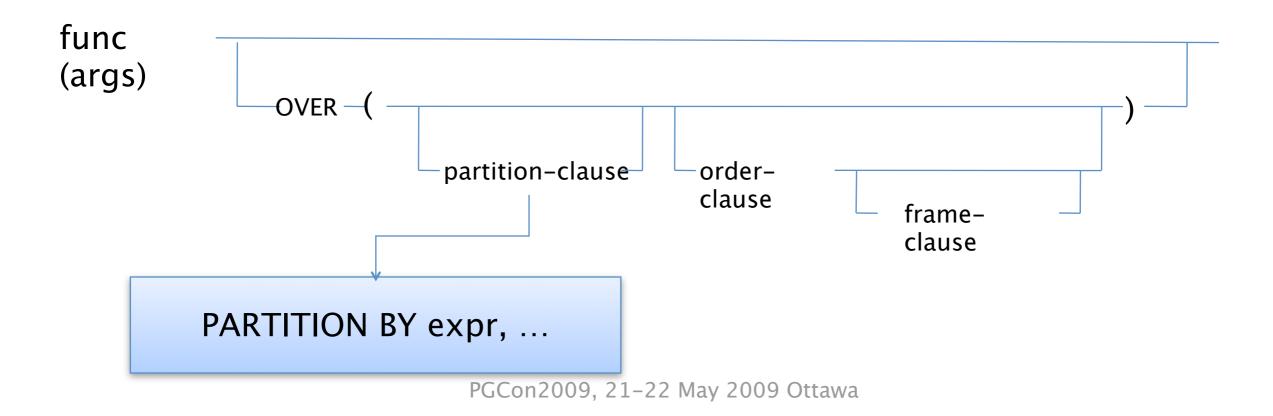
depname	empno	salary	avg	diff
develop	8	6000	5020	980
sales	6	5500	5025	475
personnel	2	3900	3700	200
develop	11	5200	5020	180
develop	10	5200	5020	180
sales	1	5000	5025	-25
personnel	5	3500	3700	-200
sales	4	4800	5025	-225
sales	3	4800	5025	-225
develop	9	4500	5020	-520
develop	7	4200	5020	-820

Anatomy of a Window

- Represents set of rows abstractly as:
- A partition
 - Specified by PARTITION BY clause
 - Never moves
 - Can contain:
- A frame
 - Specified by ORDER BY clause and frame clause
 - Defined in a partition
 - Moves within a partition
 - Never goes across two partitions
- Some functions take values from a partition.
 Others take them from a frame.

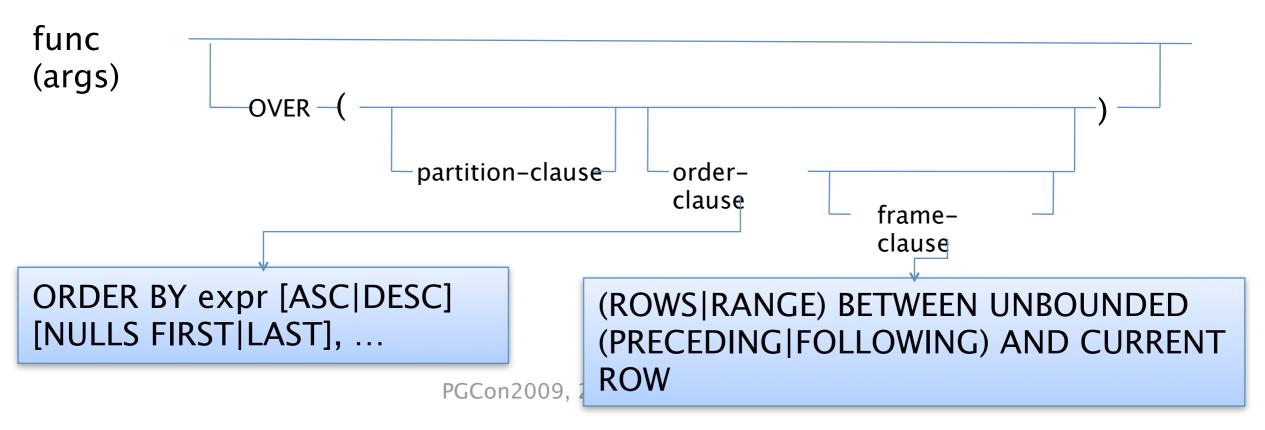
A partition

- Specified by PARTITION BY clause in OVER()
- Allows to subdivide the table, much like GROUP BY clause
- Without a PARTITION BY clause, the whole table is in a single partition



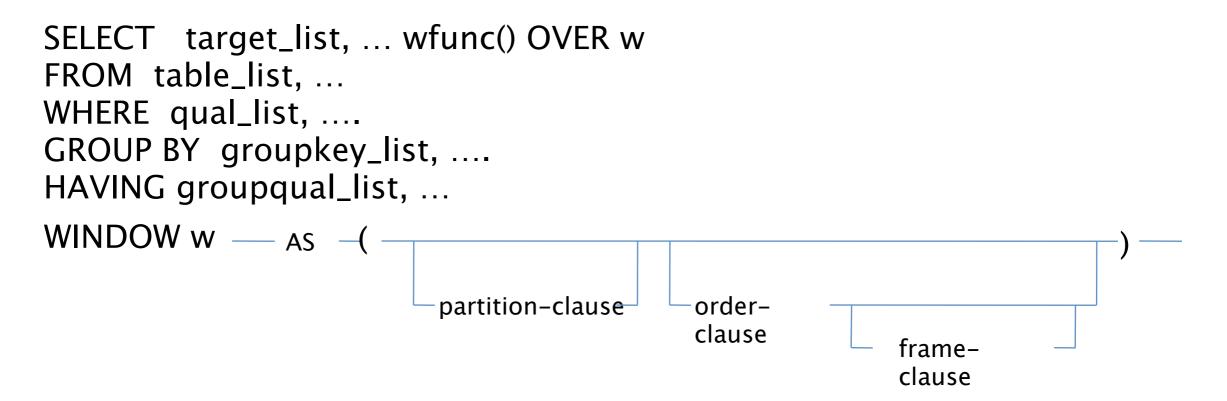
A frame

- Specified by ORDER BY clause and frame clause in OVER()
- Allows to tell how far the set is applied
- Also defines ordering of the set
- Without order and frame clauses, the whole of partition is a single frame



The WINDOW clause

- You can specify common window definitions in one place and name it, for convenience
- You can use the window at the same query level



Built-in Windowing Functions

List of built-in Windowing Functions

- row_number()
- rank()
- dense_rank()
- percent_rank()
- cume_dist()
- ntile()

- lag()
- lead()
- first_value()
- last_value()
- nth_value()

row_number()

• Returns number of the current row

SELECT val, row_number() OVER (ORDER BY val DESC) FROM tbl;

val	row_number()
5	1
5	2
3	3
1	4

Note: row_number() always incremented values independent of frame

rank()

• Returns rank of the current row with gap

SELECT val, rank() OVER (ORDER BY val DESC) FROM tbl;

val	rank()	
5	1	
5	1	gap
3	3	
1	4	

Note: rank() OVER(*empty*) returns 1 for all rows, since all rows are peers to each other

dense_rank()

• Returns rank of the current row without gap

SELECT val, dense_rank() OVER (ORDER BY val DESC) FROM tbl;

val	dense_rank()	
5	1	
5	1	no gap
3	2	
1	3	

Note: dense_rank() OVER(*empty*) returns 1 for all rows, since all rows are peers to each other

percent_rank()

Returns relative rank; (rank() – 1) / (total row – 1)

SELECT val, percent_rank() OVER (ORDER BY val DESC) FROM tbl;

val	percent_rank()
5	0
5	0
3	0.6666666666666666
1	1

values are always between 0 and 1 inclusive.

Note: percent_rank() OVER(*empty*) returns 0 for all rows, since all rows are peers to each other

cume_dist()

• Returns relative rank; (# of preced. or peers) / (total row)

SELECT val, cume_dist() OVER (ORDER BY val DESC) FROM tbl;

val	cume_dist()	
5	0.5	= 2 / 4
5	0.5	= 2 / 4
3	0.75	= 3 / 4
1	1	= 4 / 4

The result can be emulated by

"count(*) OVER (ORDER BY val DESC) / count(*) OVER ()"

Note: cume_dist() OVER(*empty*) returns 1 for all rows, since all rows are peers to each other

ntile()

• Returns dividing bucket number

SELECT val, ntile(3) OVER (ORDER BY val DESC) FROM tbl;

val	ntile(3)	
5	1	4 % 3 = 1
5	1	
3	2	
1	3	

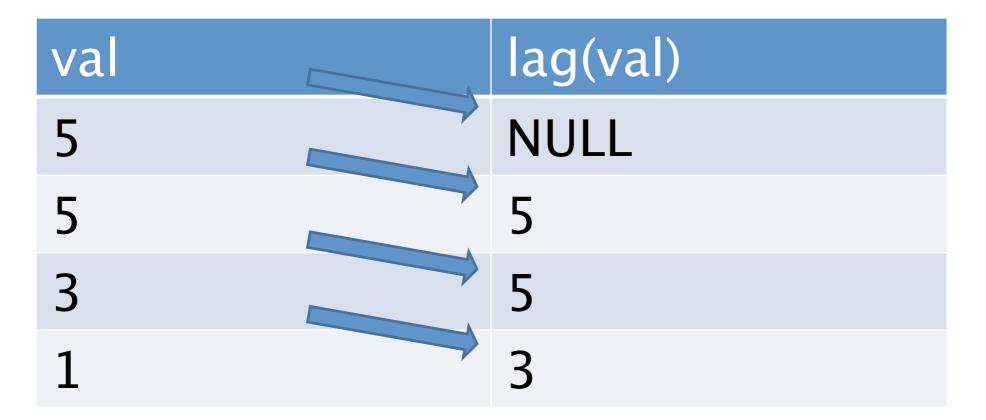
The results are the divided positions, but if there's remainder add row from the head

Note: ntile() OVER (*empty*) returns same values as above, since ntile() doesn't care the frame but works against the partition

lag()

• Returns value of row above

SELECT val, lag(val) OVER (ORDER BY val DESC) FROM tbl;

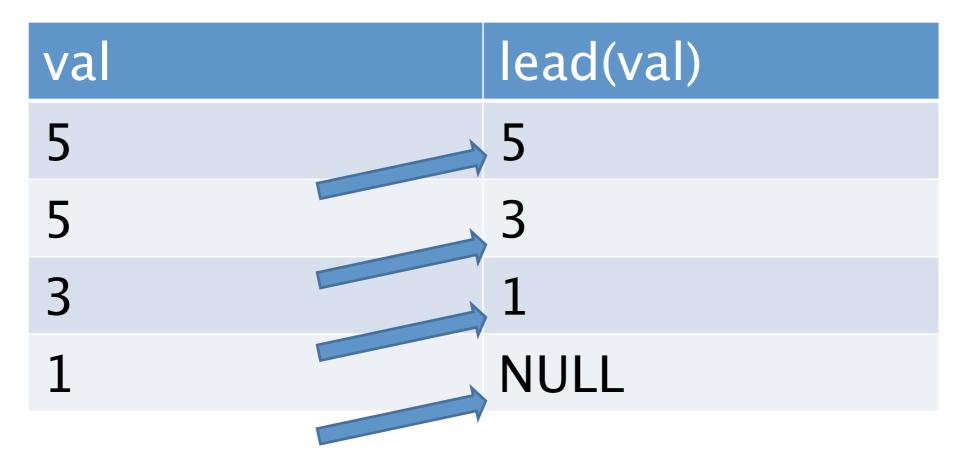


Note: lag() only acts on a partition.

lead()

• Returns value of the row below

SELECT val, lead(val) OVER (ORDER BY val DESC) FROM tbl;

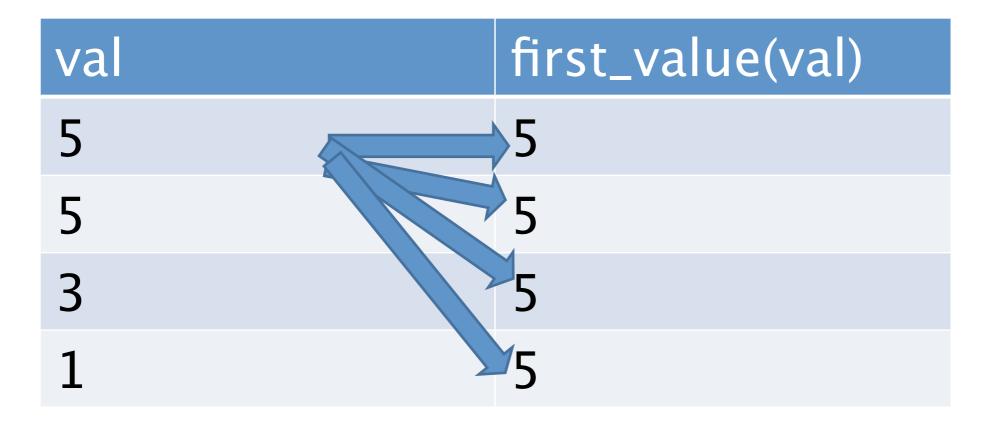


Note: lead() acts against a partition.

first_value()

• Returns the first value of the frame

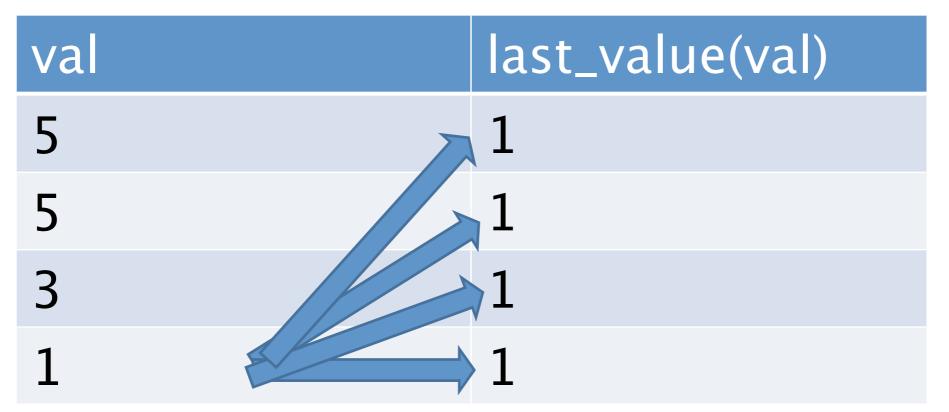
SELECT val, first_value(val) OVER (ORDER BY val DESC) FROM tbl;



last_value()

• Returns the last value of the frame

SELECT val, last_value(val) OVER (ORDER BY val DESC ROWS BETWEEN UNBOUNDED PRECEEDING AND UNBOUNDED FOLLOWING) FROM tbl;

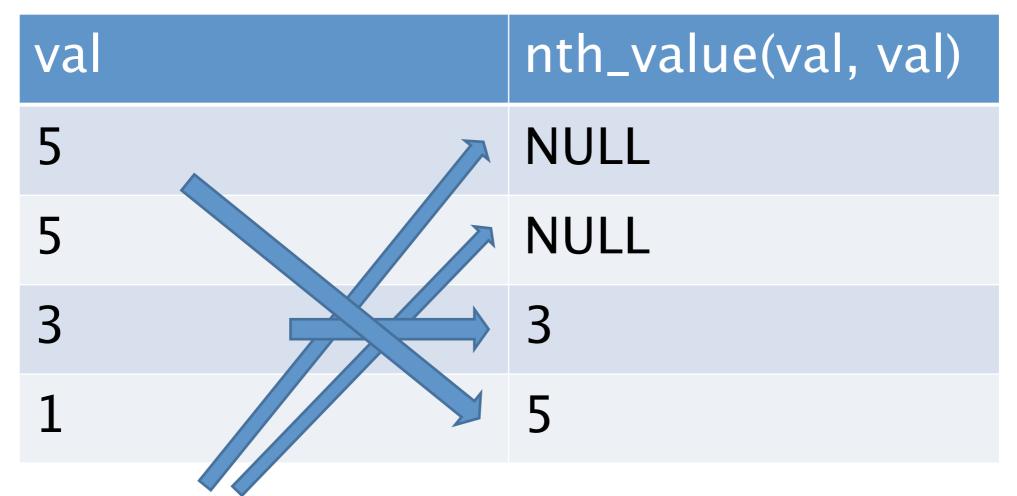


Note: frame clause is necessary since you have a frame between the first row and the current row by only the order clause

nth_value()

• Returns the n-th value of the frame

SELECT val, nth_value(val, val) OVER (ORDER BY val DESC ROWS BETWEEN UNBOUNDED PRECEEDING AND UNBOUNDED FOLLOWING) FROM tbl;



Note: frame clause is necessary since you have a frame between the first row and the current row by only the order clause PGCon2009, 21–22 May 2009 Ottawa

aggregates(all peers)

• Returns the same values along the frame

SELECT val, sum(val) OVER () FROM tbl;

val	sum(val)
5	14
5	14
3	14
1	14

Note: all rows are the peers to each other

cumulative aggregates

• Returns different values along the frame

SELECT val, sum(val) OVER (ORDER BY val DESC) FROM tbl;

val	sum(val)
5	10
5	10
3	13
1	14

Note: row#1 and row#2 return the same value since they are the peers. the result of row#3 is sum(val of row#1...#3)

Inside the Implementation

Where are they?

- Windowing Functions are identified by a flag in pg_proc, which means they are very similar to plain functions
 - pg_proc.proiswindow : boolean
- All the aggregates including user-defined ones in any language can be called in WindowAgg, too
 - Your old code gets new behaviour!
- src/include/catalog/pg_proc.h
 - pg_proc catalog definitions
- src/backend/utils/adt/windowfuncs.c
 - implementations of built-in window functions

Added relevant nodes

• WindowFunc

- primitive node for function itself "wfunc(avg1, ...)"

• WindowDef

- parse node for window definition "over (partition by ... order by ...)"
- WindowClause
 - parse node for window clause "window (partition by ... order by ...) as w"
- WindowAgg
 - plan node for window aggregate
- WindowAggState
 - executor node for window aggregate

Hacking the planner

SELECT depname, empno, salary, avg(salary) over (partition by depname) AS a, rank() over (partition by depname order by salary desc) AS r FROM empsalary ORDER BY r

Sort (cost=215.75..218.35 rows=1040 width=44)
Output: depname, empno, salary, (avg(salary) OVER (?)), (rank() OVER (?))
Sort Key: (rank() OVER (?))
-> WindowAgg (cost=142.83..163.63 rows=1040 width=44)
Output: depname, empno, salary, (avg(salary) OVER (?)), rank() OVER (?)
-> Sort (cost=142.83..145.43 rows=1040 width=44)
Output: depname, empno, salary, (avg(salary) OVER (?))
Sort Key: depname, salary
-> WindowAgg (cost=72.52..90.72 rows=1040 width=44)
Output: depname, empno, salary, avg(salary) OVER (?)
-> Sort (cost=72.52..75.12 rows=1040 width=44)
Output: depname, empno, salary
-> Sort (cost=72.52..75.12 rows=1040 width=44)
Output: depname, empno, salary
Sort Key: depname
-> Seq Scan on empsalary (cost=0.00..20.40 rows=1040 width=44)
Output: depname, empno, salary

Hacking the planner

Final output

TargetEntry1:depname

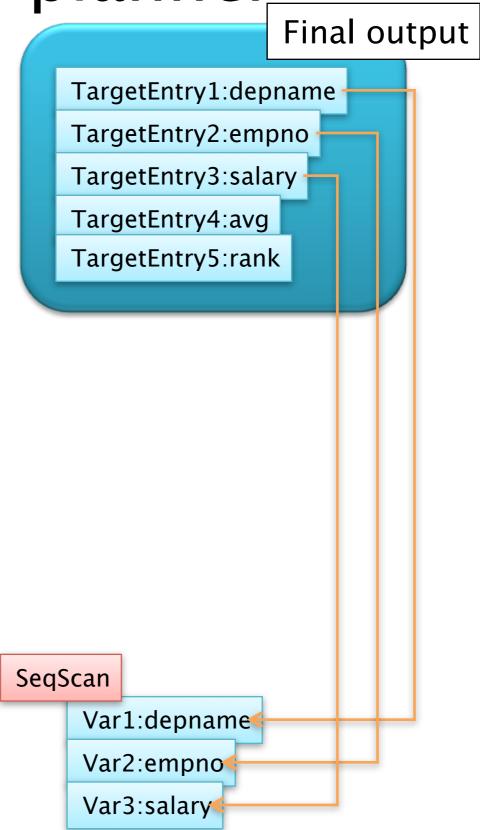
TargetEntry2:empno

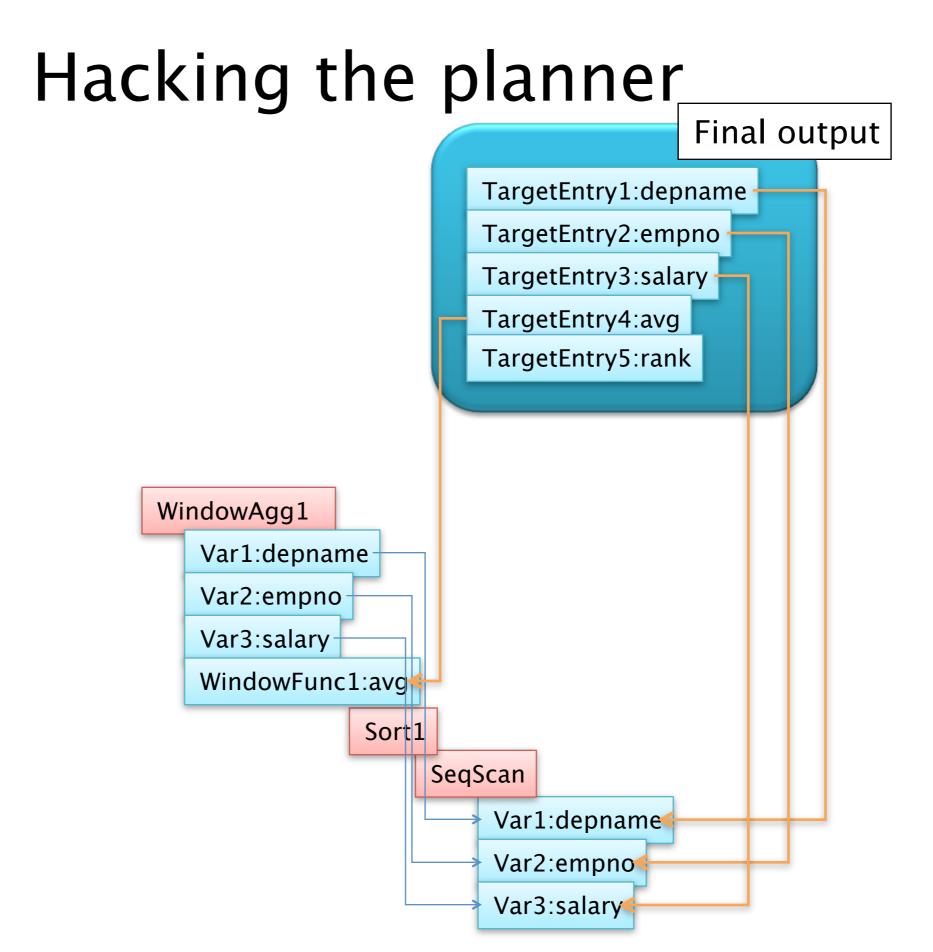
TargetEntry3:salary

TargetEntry4:avg

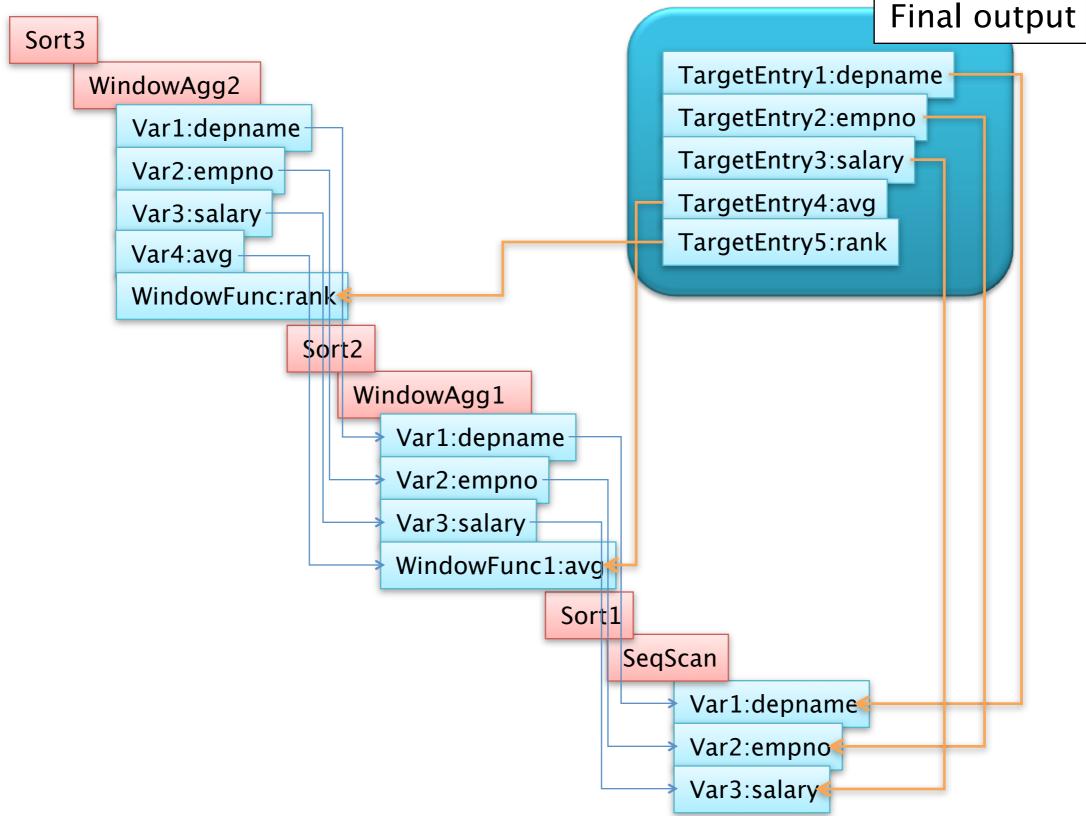
TargetEntry5:rank

Hacking the planner





Hacking the planner



Hacking the planner

• We could optimize it by relocating WindowAgg

```
SELECT depname, empno, salary,
rank() over (partition by depname order by salary desc) AS r,
avg(salary) over (partition by depname) AS a
FROM empsalary ORDER BY r
```

```
Sort (cost=161.03..163.63 rows=1040 width=44)

Output: depname, empno, salary, (rank() OVER (?)), (avg(salary) OVER (?))

Sort Key: (rank() OVER (?))

-> WindowAgg (cost=72.52..108.92 rows=1040 width=44)

Output: depname, empno, salary, (rank() OVER (?)), avg(salary) OVER (?) NO Sort!

-> WindowAgg (cost=72.52..93.32 rows=1040 width=44)

Output: depname, empno, salary, rank() OVER (?)

-> Sort (cost=72.52..75.12 rows=1040 width=44)

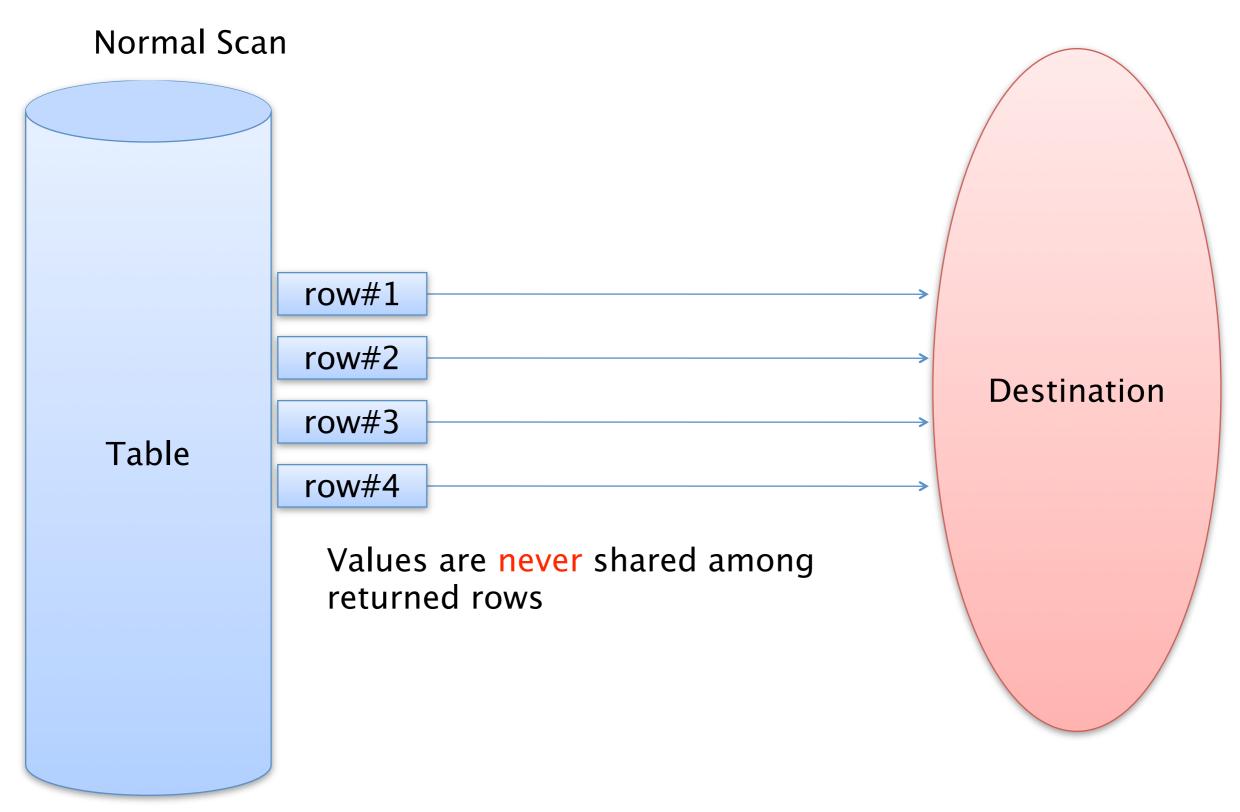
Output: depname, empno, salary

Sort Key: depname, salary

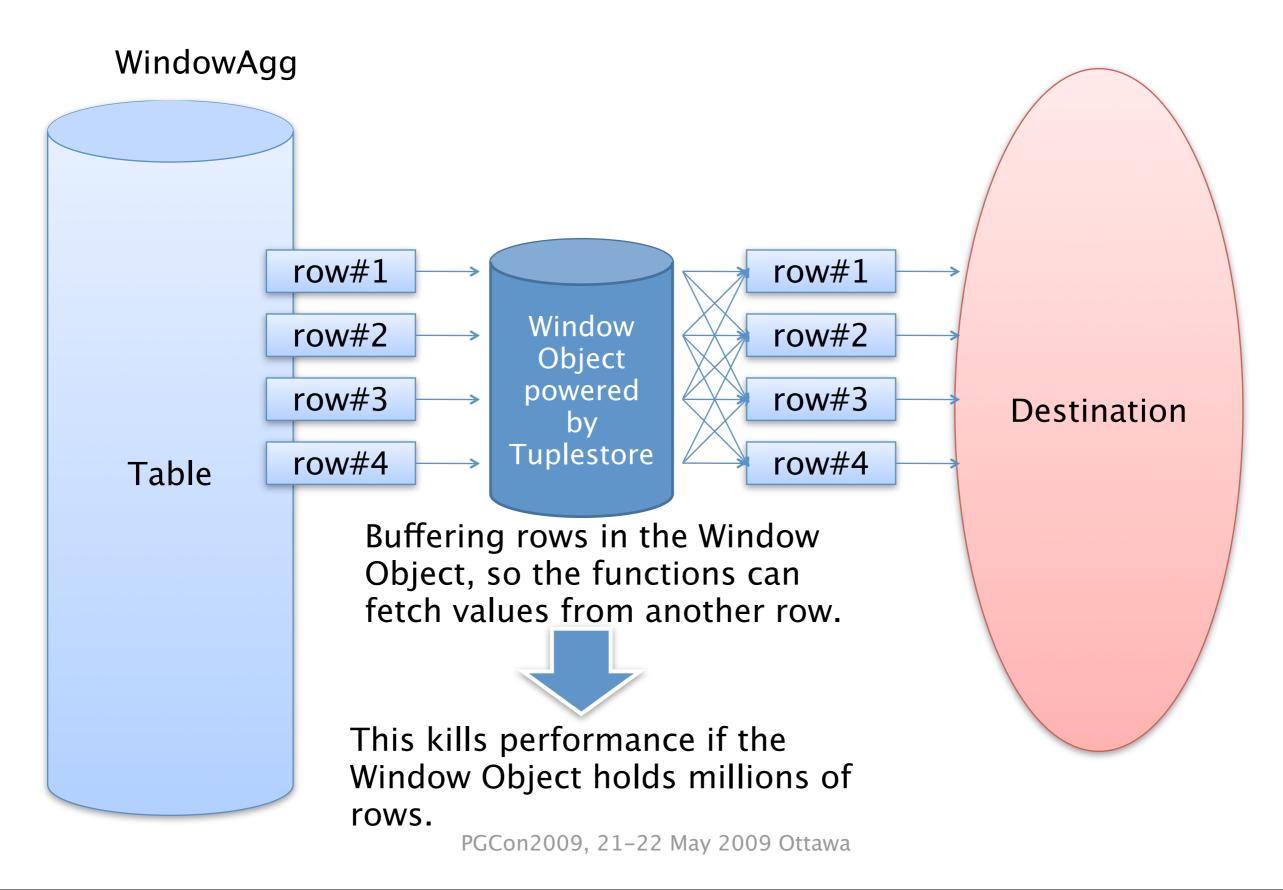
-> Seq Scan on empsalary (cost=0.00..20.40 rows=1040 width=44)

Output: depname, empno, salary
```

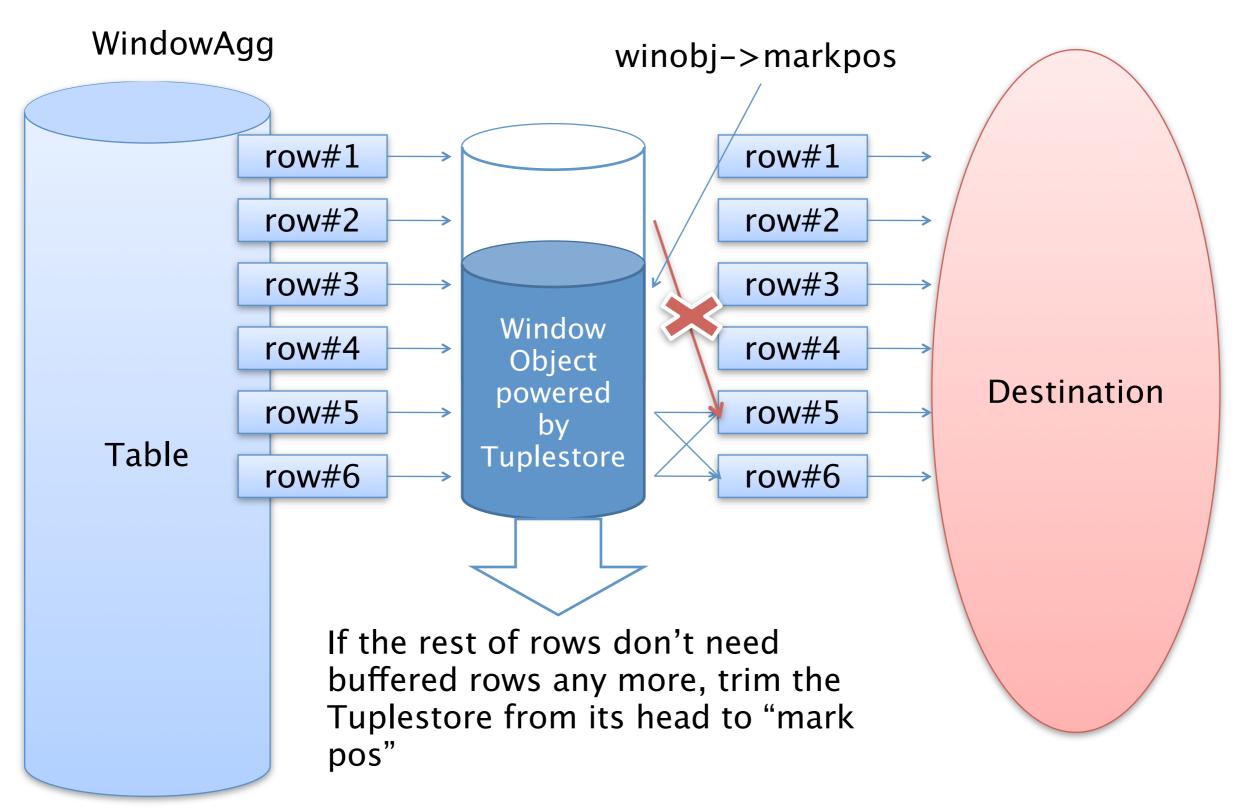
How the executor creates a window



How the executor creates a window

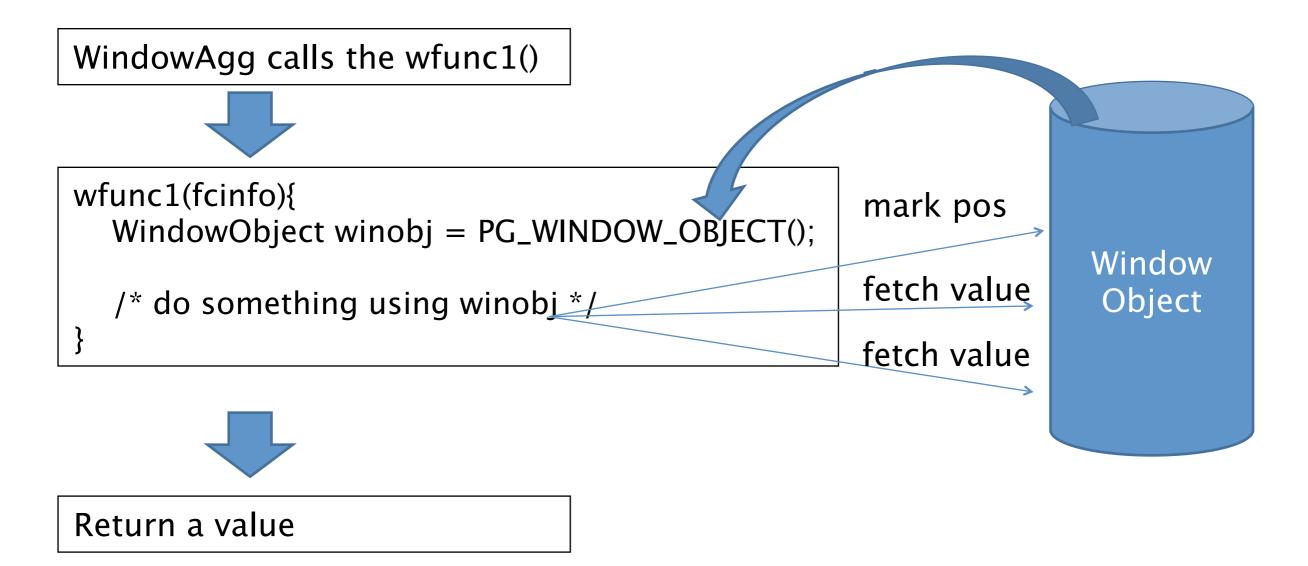


How the executor creates a window



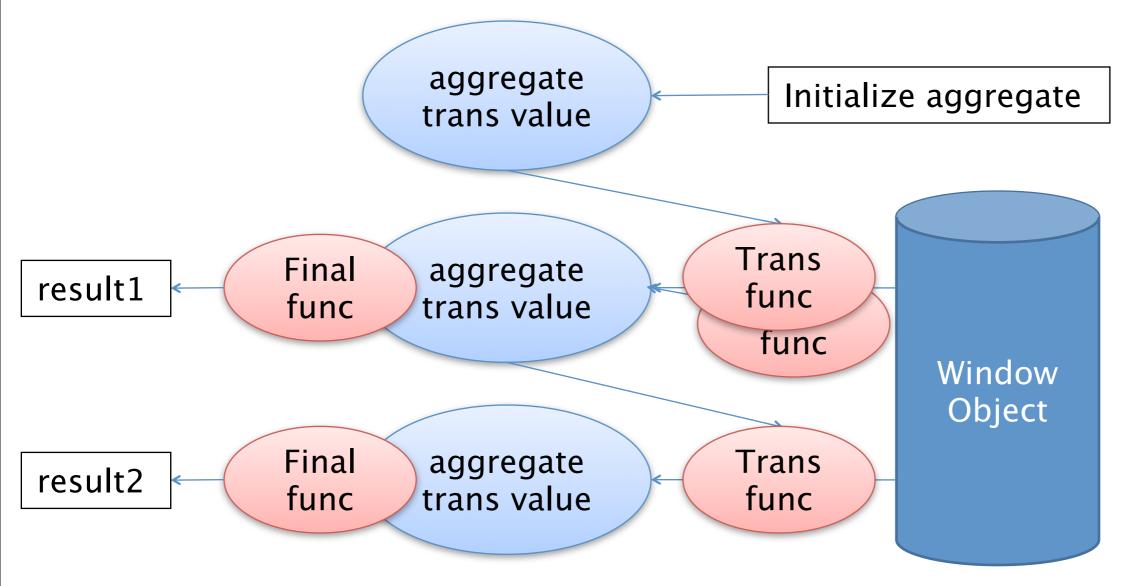
How the Windowing Functions work

- Functions workflow is very similar to plain functions
- Each function has its own seek position



How the Windowing Functions work

• What about aggregates?



Trans value is cached and reused after final function call.

-> BE CAREFUL of anything that will break if the final function is called more than once.

http://archives.postgresgl.org//pgsgl-hackers/2008-12/msg01820.php

Write Your Own Windowing Function

Write your own Windowing Function

- NOT documented for now, but you can do it
- The normal argument API like PG_GETARG_XXX() does not work for window functions
- STRICT-ness is meaningless
- The function must use the V1 calling convention
- Get WindowObject using PG_WINDOW_OBJECT()
- Check window-ness by WindowObjectIsValid()
- Call APIs to fetch values
- More details in src/include/windowapi.h, see also:
 - src/executor/nodeWindowAgg.c
 - src/utils/adt/windowfuncs.c
- These may be changed in the future releases

Windowing Function APIs

- PG_WINDOW_OBJECT(fcinfo)
 - Retrieve WindowObject, which is an interface of the window, for this call.
- WinGetPartitionLocalMemory(winobj, sz)
 - Store its own memory. It is used in rank() to save the current rank, for example.
- WinGetCurrentPosition(winobj)
 - Get current position in the partition (not in the frame).
 Same as row_number()
- WinGetPartitionRowCount(winobj)
 - Count up how many rows in the partition. Used in ntile() for example.

Windowing Function APIs

- WinSetMarkPosition(winobj, markpos)
 - Give to winobj a hint that you don't want rows preceding to markpos anymore
- WinRowsArePeers(winobj, pos1, pos2)
 - Rows at pos1 and at pos2 are peers? "Peers" means both rows have the same value in the meaning of ORDER BY columns

Windowing Function APIs

- WinGetFuncArgInPartition(winobj, argno, relpos, seektype, set_mark, isnull, isout)
 - Fetch argument from another row in the partition. "isout" will be set to true if "relpos" is out of range of the partition.
- WinGetFuncArgInFrame(winobj, argno, relpos, seektype, set_mark, isnull, isout)
 - Fetch argument from another row in the frame. "isout" will be set to true if "relpos" is out of range of the frame (may be within the partition.)
- WinGetFuncArgCurrent(winobj, argno, isnull)
 - Same as PG_GETARG_XXX(argno). Fetch argument from the current row.

Moving average example

- <u>http://archives.postgresql.org/pgsql-</u> <u>hackers/2009-01/msg00562.php</u>
- With standard SQL, you have to wait for extended frame clause support like ROWS n (PRECEDING|FOLLOWING) to calculate curve smoothing, but this sample does it now.

Future work

Future work for later versions

- More support for frame clause
 - ROWS n (PRECEDING|FOLLOWING)
 - RANGE val (PRECEDING|FOLLOWING)
 - EXCLUDE (CURRENT ROW|GROUP|TIES|NO OTHERS)
 - Not so difficult for window functions but integration with classical aggregates is hard
- Performance improvements
 - Reduce tuplestore overhead
 - Relocate WindowAgg node in the plan
- Support PLs for writing Windowing Functions
 - Let PLs call Windowing Function APIs

Finally...

Thanks all the hackers!!

• Simon Riggs

discussion over my first ugly and poor design

• Heikki Linnakangas

- improvement of window object and its buffering strategy
- Tom Lane
 - code rework overall and deep reading in the SQL spec
- David Fetter
 - help miscellaneous things like git repository and this session
- David Rowley
 - many tests to catch early bugs, as well as spec investigations
- Takahiro Itagaki, Pavel Stehule, and All the Hackers
- I might have missed your name here, but really appreciate your help.